

**Name:**

**Date:**

**MA 515, Linear Programming  
and Combinatorial Optimization  
Fall 2008, Final Exam**

There are 150 possible points. You are not allowed to use your book or notes.

**Problem 1: (30 points)** First, describe the main themes of this course and their interactions. Second, choose a theorem that you feel excellently represents these themes and do the following:

- (1) State the theorem;
- (2) Give an example of the theorem in action; and
- (3) Explain why you chose this theorem, specifically addressing how it relates to the main themes you described in the first part.

**Problem 2: (40 points)** Consider the following five contexts in which we have a “max=min” result.

- (1) Linear programming
- (2) Greedy algorithm for matroids
- (3) Matroid intersection
- (4) Bipartite matchings
- (5) Matchings in general graphs

Choose *four* of these and make a precise statement of the “max=min” result in each case. Sketch the proof for *one* of these, indicating clearly which one you have chosen.

Complete four of the following five problems, each worth 20 points.

**Problem 3a:** Consider the following matrix:

$$A := \begin{bmatrix} -1 & -1 & -1 & 0 \\ 1 & 0 & 0 & -1 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 \end{bmatrix}$$

Let  $M$  be the linear matroid determined by  $A$ . Provide a description of the matroid polytope for  $M$  as an  $\mathcal{H}$ -polytope, and use linear programming to find a maximal independent set in  $M$ .

**Problem 3b:** Construct a graph  $G$  on 6 vertices with 9 edges. Assign to each of these edges a distinct positive weight. Choose a specific vertex, label it  $v$ , and use Dijkstra's algorithm to find all minimum-weight  $v - w$ -paths in  $G$ .

**Problem 3c:** Let  $G$  be a planar graph. Take any planar embedding of  $G$  and form the planar dual  $G^*$ . Prove that the graphic matroid of  $G^*$  is the dual of the graphic matroid of  $G$ .

**Problem 3d:** Let  $G$  be a graph with no loops and no multiple edges. Prove that a matching  $S$  of  $G$  is of maximum cardinality if and only if  $G$  has no augmenting path with respect to  $S$ .

**Problem 3e:** Given two matroids  $M_1$  and  $M_2$  on a common ground set  $E$  and an independent set  $S \in \mathcal{I}(M_1) \cap \mathcal{I}(M_2)$ , define the bipartite augmentation digraph for  $S$  and explain its role in the Cardinality Matroid-Intersection Algorithm.