

Lower Bounds For the Condition Number of  
a Real Confluent Vandermonde Matrix<sup>1</sup>

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ABSTRACT

Lower bounds on the condition number  $\kappa_p(V_c)$  of a real confluent Vandermonde matrix  $V_c$  are established in terms of the dimension  $n$  or  $n$  and the largest absolute value among all nodes that define the confluent Vandermonde matrix or  $n$  and the interval that contains the nodes. In particular, it is proved that for any modest  $k_{\max}$  (the largest number of equal nodes),  $\kappa_p(V_c)$  behaves no smaller than  $\mathcal{O}_n((1 + \sqrt{2})^n)$ , and than  $\mathcal{O}_n((1 + \sqrt{2})^{2n})$  if all nodes are nonnegative. It is not clear whether those bounds are asymptotically sharp for modest  $k_{\max}$ .

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<sup>1</sup>This report is available on the web at <http://www.ms.uky.edu/~math/MAreport/PDF/2004-10.pdf>.

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# Lower Bounds For the Condition Number of a Real Confluent Vandermonde Matrix

Ren-Cang Li \*

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## Abstract

Lower bounds on the condition number  $\kappa_p(V_c)$  of a real confluent Vandermonde matrix  $V_c$  are established in terms of the dimension  $n$  or  $n$  and the largest absolute value among all nodes that define the confluent Vandermonde matrix or  $n$  and the interval that contains the nodes. In particular, it is proved that for any modest  $k_{\max}$  (the largest number of equal nodes),  $\kappa_p(V_c)$  behaves no smaller than  $\mathcal{O}_n((1+\sqrt{2})^n)$ , and than  $\mathcal{O}_n((1+\sqrt{2})^{2n})$  if all nodes are nonnegative. It is not clear whether those bounds are asymptotically sharp for modest  $k_{\max}$ .

## 1 Introduction

Given  $n$  numbers  $\alpha_1, \alpha_2, \dots, \alpha_n$  called *nodes*, the associated *Vandermonde Matrix* is defined as

$$V \stackrel{\text{def}}{=} \begin{pmatrix} 1 & 1 & \cdots & 1 \\ \alpha_1 & \alpha_2 & \cdots & \alpha_n \\ \vdots & \vdots & \ddots & \vdots \\ \alpha_1^{n-1} & \alpha_2^{n-1} & \cdots & \alpha_n^{n-1} \end{pmatrix}. \quad (1.1)$$

It, for example, arises from polynomial interpolation and others [4].  $V$  is invertible if all nodes  $\alpha_j$  are distinct, i.e.,  $\alpha_i \neq \alpha_j$  for  $i \neq j$ , but it becomes singular whenever  $\alpha_i = \alpha_j$  for some  $i \neq j$ . A generalization of  $V$  for nodes not all of which are distinct is the so-called *Confluent Vandermonde Matrices*, e.g.,

$$\begin{pmatrix} 1 & 0 & 0 & 1 & 1 & 0 \\ \alpha_1 & 1 & 0 & \alpha_4 & \alpha_5 & 1 \\ \alpha_1^2 & 2\alpha_1 & 2 & \alpha_4^2 & \alpha_5^2 & 2\alpha_5 \\ \alpha_1^3 & 3\alpha_1^2 & 6\alpha_1 & \alpha_4^3 & \alpha_5^3 & 3\alpha_5^2 \\ \alpha_1^4 & 4\alpha_1^3 & 12\alpha_1^2 & \alpha_4^4 & \alpha_5^4 & 4\alpha_5^3 \\ \alpha_1^5 & 5\alpha_1^4 & 20\alpha_1^3 & \alpha_4^5 & \alpha_5^5 & 5\alpha_5^4 \end{pmatrix},$$

where  $\alpha_1 = \alpha_2 = \alpha_3$  and  $\alpha_5 = \alpha_6$ . The second, third, and sixth columns are obtained by “differentiating” the previous column. Confluent Vandermonde matrices arises in Hermite

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interpolation [3], for example. Adopting the formulation in [8], we define *confluent Vandermonde matrix*  $V_c$  as follows. First

$$\boxed{\begin{aligned} \{\alpha_j\}_{j=1}^n \text{ are ordered so that equal nodes are contiguous, i.e.,} \\ \alpha_i = \alpha_j \quad (i < j) \quad \Rightarrow \quad \alpha_i = \alpha_{i+1} = \cdots = \alpha_j. \end{aligned}} \quad (1.2)$$

Define

$$V_c = (f_1(\alpha_1) \ f_2(\alpha_2) \ \cdots \ f_n(\alpha_n)), \quad (1.3)$$

where vector function  $f_j(t)$  is defined recursively by

$$f_j(t) = \begin{cases} (1 \ t \ \cdots \ t^{n-1})^T, & \text{if } j = 1 \text{ or } \alpha_j \neq \alpha_{j-1}, \\ \frac{d}{dx} f_{j-1}(t), & \text{otherwise,} \end{cases} \quad (1.4)$$

where “ $\cdot^T$ ” is the transpose of a vector or matrix. As far as defining  $V_c$  is concerned,  $\alpha_j$  can be real or complex. But in this paper, we shall focus on real  $\alpha_j$ . In what follows,  $\alpha_j$  and  $V_c$ , as well as  $\alpha \stackrel{\text{def}}{=} \max_j |\alpha_j|$ , are reserved for their assignments here.

(Optimal) condition numbers for real Vandermonde matrices have been systematically studied by Gautschi and his co-author (see [7] and references therein), and more recently by Tyrtyshnikov [12], Beckermann [2], and Li [10]. In this paper, we shall establish several lower bounds on the  $\ell_p$ -condition number  $\kappa_p(V_c) \equiv \|V_c\|_p \|V_c^{-1}\|_p$  in terms of  $n$  or  $n$  and  $\alpha$  or  $n$  and the interval  $[a, b]$  that contains all nodes. In particular, we will show that for fixed  $k_{\max}$  (the largest number of equal nodes),  $\kappa_p(V_c)$  behaves no smaller than  $\mathcal{O}_n((1 + \sqrt{2})^n)$ , where notation  $a_n = \mathcal{O}_n(b_n)$  means  $c_1 n^{d_1} \leq a_n/b_n \leq c_2 n^{d_2}$  for some constants  $c_1, c_2, d_1$ , and  $d_2$ . We also obtain a qualitative plot in Figure 1.1 to show how our lower bounds on  $\min_{\alpha_j} \kappa_p(V_c)$  and  $\min_{\alpha_j \geq 0} \kappa_p(V_c)$  behave qualitatively as functions of  $\alpha$ . What Figure 1.1 says that initially as  $\alpha$  increases, our lower bound for  $\min_{\alpha_j} \kappa_p(V_c)$  and that for  $\min_{\alpha_j \geq 0} \kappa_p(V_c)$  decrease until at  $\alpha = \alpha_{\text{opt}}$  when global minimums of the bounds are reached, and then they start climbing again. Notice  $\alpha_{\text{opt}}$  may be different for the two cases, but  $\alpha_{\text{opt}} = \mathcal{O}(1)$  in both cases. What that is not clear, however, is how sharp our lower bounds are, in contrast to many of those in [10] for Vandermonde matrices that were proved to be asymptotically optimal.

Optimally conditioned confluent Vandermonde matrices can be much worse ill-conditioned than optimally conditioned Vandermonde matrices. One extreme example would be that all nodes are equal  $\alpha_1 = \cdots = \alpha_n$  for which  $V_c$  is lower triangular and thus

$$\kappa_p(V_c) \geq (n-1)! \sim \sqrt{2\pi} n^{n-1/2} e^{-n}$$

by Stirling’s asymptotic formula [1, Page 18], and it becomes an equality for  $\alpha_1 = \cdots = \alpha_n = 0$ . While for optimally conditioned Vandermonde matrices,  $\kappa_p(V)$  goes to  $\infty$  as fast as  $(1 + \sqrt{2})^n$  modulo a factor  $n^d$  for  $d \leq 1$  [2, 10].

The rest of this paper is organized as follows. Section 2 reviews some preliminary results from [10] in connection to the coefficients of translated Chebyshev polynomials. Two general lower bounds on  $\kappa_p(V_c)$  are established in Section 3, but they are not uniform. Uniform bounds for  $p = \infty$  are obtained in Section 4 for all real  $V_c$  and in Section 5 for  $V_c$  with nonnegative nodes.

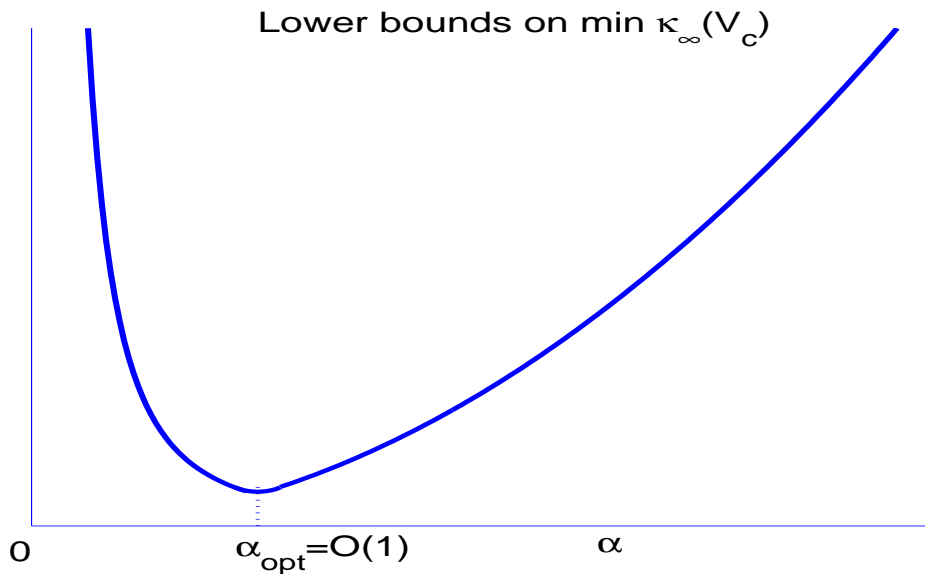


Figure 1.1: Qualitative behaviors of our lower bounds on  $\min_{\alpha_j} \kappa_p(V_c)$  and  $\min_{\alpha_j \geq 0} \kappa_p(V_c)$  (subject to fixed  $\alpha$ ) as  $\alpha$  varies for any given  $k_{\max}$ .

## 2 Preliminaries

Let us start by briefly reviewing relevant notation and results from [10]. Let  $[a, b]$  be the interval in which lie all  $\alpha_j$ .  $T_n(t) = \cos(n \arccos t)$  is the  $n$ th Chebyshev polynomial of the first kind. Define the  $n$ th translated Chebyshev polynomial  $T_n(x; \omega, \tau) \stackrel{\text{def}}{=} T_n(x/\omega + \tau)$ , where

$$\omega = \frac{b-a}{2}, \quad \tau = -\frac{b+a}{b-a}.$$

Let  $a_{jn} \equiv a_{jn}(\omega, \tau)$  be the coefficient of  $x^j$  in  $T_n(x; \omega, \tau)$ , i.e.,

$$T_n(x; \omega, \tau) = a_{nn}x^n + a_{n-1n}x^{n-1} + \cdots + a_{1n}x + a_{0n}, \quad (2.1)$$

Let  $1 \leq p \leq \infty$ . In [10],  $S_{n,p}(\omega, \tau)$  is defined as

$$S_{n,p}(\omega, \tau) = \left( \sum_{j=0}^n |a_{jn}|^p \right)^{1/p}.$$

Also explicit formulas were found for  $p = 1$  and  $-a = b$  (for which  $\omega = 0$ ):

$$S_{n,1}(\omega, 0) = \frac{1}{2} \left[ \left( \frac{1}{|\omega|} + \sqrt{1 + \frac{1}{|\omega|^2}} \right)^n + (-1)^n \left( \frac{1}{|\omega|} + \sqrt{1 + \frac{1}{|\omega|^2}} \right)^{-n} \right], \quad (2.2)$$

$$\sim \frac{1}{2} \left( \frac{1}{|\omega|} + \sqrt{1 + \frac{1}{|\omega|^2}} \right)^n. \quad (2.3)$$

and for  $a \geq 0$  (for which  $\tau \leq -1$ ):

$$S_{n,1}(\omega, \tau) = \frac{1}{2} \left[ \left( \frac{1}{|\omega|} + |\tau| \right) + \sqrt{\left( \frac{1}{|\omega|} + |\tau| \right)^2 - 1} \right]^n + \frac{1}{2} \left[ \left( \frac{1}{|\omega|} + |\tau| \right) - \sqrt{\left( \frac{1}{|\omega|} + |\tau| \right)^2 - 1} \right]^{-n}, \quad (2.4)$$

$$\sim \frac{1}{2} \left[ \left( \frac{1}{|\omega|} + |\tau| \right) + \sqrt{\left( \frac{1}{|\omega|} + |\tau| \right)^2 - 1} \right]^n. \quad (2.5)$$

For any other  $p$ ,  $S_{n,p}(\omega, \tau)$  relates to  $S_{n,1}(\omega, \tau)$  by

$$(n+1)^{-1/p'} S_{n,1}(\omega, \tau) \leq S_{n,p}(\omega, \tau) \leq S_{n,1}(\omega, \tau), \quad (2.6)$$

$$\lceil (n+1)/2 \rceil^{-1/p'} S_{n,1}(\omega, 0) \leq S_{n,p}(\omega, 0) \leq S_{n,1}(\omega, 0), \quad (2.7)$$

where  $1/p + 1/p' = 1$ , and  $\lceil \xi \rceil$  is the smallest integer that is larger than  $\xi$ .

The  $\ell_p$ -norm of vector  $u = (\mu_1 \ \mu_2 \ \cdots \ \mu_n)^T$  is defined as

$$\|u\|_p = \left( \sum_{j=1}^n |\mu_j|^p \right)^{1/p},$$

and  $\|u\|_\infty = \lim_{p \rightarrow \infty} \|u\|_p = \max_j |\mu_j|$ . The associated  $\ell_p$ -operator norm of  $m \times n$  matrix  $A$  is defined as

$$\|A\|_p = \max_{u \neq 0} \frac{\|Au\|_p}{\|u\|_p}. \quad (2.8)$$

It can be proved that [9]  $\|A\|_p = \|A^T\|_{p'}$ , and [10]

$$n^{-2/p} \kappa_p(V_c) \leq \kappa_\infty(V_c) \leq n^{2/p} \kappa_p(V_c) \quad (2.9)$$

which is useful in deriving bounds on  $\kappa_p(V_c)$  from these for  $\kappa_\infty(V_c)$  as in Sections 4 and 5.

### 3 Lower bounds on $\kappa_p(V_c)$

For the sake of presentation, we assume, in addition to (1.2),

There are  $\ell$  distinct nodes  $\alpha_j$ , having multiplicities  $k_1, k_2, \dots, k_\ell$ , respectively, where  $k_1 + k_2 + \dots + k_\ell = n$ . This implies that the first  $k_1$   $\alpha_j$ 's are equal, the next  $k_2$   $\alpha_j$ 's are equal, and so on. Define  $k_{\max} = \max_j k_j$ .

(3.1)

**Lemma 3.1** *Assume (1.2) and (3.1). Then*

$$\|V_c\|_p \geq \max \left\{ \ell^{1/p'}, \alpha^{n-1} \right\}, \quad (3.2)$$

$$\|V_c\|_p \geq \left( \sum_{j=1}^n \alpha^{(j-1)p} \right)^{1/p}. \quad (3.3)$$

*Proof:* Let  $e_j$  be the  $j$ th column of the  $n \times n$  identity matrix  $I_n$  (or simply  $I$  if  $n$  is clear from the context). Use  $\|V_c\|_p \geq \|V_c^T e_1\|_{p'}$  and  $\|V_c\|_p \geq \|V_c^T e_n\|_{p'}$  to get (3.2), and use  $\|V_c\|_p \geq \max_j \|V_c^T e_j\|_{p'}$  to get (3.3).  $\blacksquare$

**Lemma 3.2** For  $0 \leq k \leq n$ ,

$$\left| \frac{d}{dx^k} T_n(x; \omega, \tau) \right| \leq \frac{[n(n-1) \cdots (n-k+1)]^2}{\omega^k} \quad \text{for } x \in [a, b]. \quad (3.4)$$

*Proof:* It follows from  $T_n(x; \omega, \tau) = T_n(x/\omega + \tau) \equiv T_n(t)$  that

$$\frac{d^k}{dx^k} T_n(x; \omega, \tau) = \frac{1}{\omega^k} T_n^{(k)}(t),$$

where  $t \equiv t(x) = x/\omega + \tau$ . It suffices to show that  $|T_n^{(k)}(t)| \leq [n(n-1) \cdots (n-k+1)]^2$  for  $t \in [-1, 1]$  since  $t(x)$  maps  $x \in [a, b]$  to  $t \in [-1, 1]$ . By Markov's Inequality [5, Page 233],

$$\begin{aligned} \max_{t \in [-1, 1]} |T_n^{(k)}(t)| &\leq (n-k+1)^2 \max_{t \in [-1, 1]} |T_n^{(k-1)}(t)| \\ &\leq \cdots \\ &\leq [n(n-1) \cdots (n-k+1)]^2 \max_{t \in [-1, 1]} |T_n(t)| \\ &= [n(n-1) \cdots (n-k+1)]^2, \end{aligned}$$

as expected.  $\blacksquare$

**Lemma 3.3** Assume (1.2) and (3.1). Then

$$\|V_c^{-1}\|_p \geq \min_{1 \leq k \leq k_{\max}} \left[ \frac{(n-k)!}{(n-1)!} \right]^2 \omega^{k-1} \times \frac{S_{n-1, p'}(\omega, \tau)}{n^{1/p'}}. \quad (3.5)$$

*Proof:* Let  $v$  be the vector of the coefficients of the translated Chebyshev polynomial  $T_{n-1}(x; \omega, \tau)$ , i.e.,  $v = (a_{0n-1} \ a_{1n-1} \ \cdots \ a_{n-1n-1})^T$ . Then

$$V_c^T v = (T_{n-1}(\alpha_1; \omega, \tau) \ T'_{n-1}(\alpha_1; \omega) \ \cdots \ T_{n-1}^{(k_1-1)}(\alpha_1; \omega) \ \cdots \ \cdots)^T$$

which yields, by Lemma 3.2, for  $1 \leq p' < \infty$

$$\|V_c^T v\|_{p'}^{p'} \leq \sum_{j=1}^{\ell} \left( 1^{p'} + \left[ \frac{(n-1)^2}{\omega} \right]^{p'} + \cdots + \left[ \frac{[(n-1)(n-2) \cdots (n-k_j+1)]^2}{\omega^{k_j-1}} \right]^{p'} \right) \quad (3.6)$$

$$\begin{aligned} &\leq \sum_{j=1}^{\ell} k_j \times \left( \max_{1 \leq k \leq k_j} \left[ \frac{(n-1)!}{(n-k)!} \right]^2 \frac{1}{\omega^{k-1}} \right)^{p'} \\ &\leq n \times \left( \max_{1 \leq k \leq k_{\max}} \left[ \frac{(n-1)!}{(n-k)!} \right]^2 \frac{1}{\omega^{k-1}} \right)^{p'} \end{aligned} \quad (3.7)$$

which gives

$$\|V_c^T v\|_{p'} \leq n^{1/p'} \times \max_{1 \leq k \leq k_{\max}} \left[ \frac{(n-1)!}{(n-k)!} \right]^2 \frac{1}{\omega^{k-1}}. \quad (3.8)$$

This is proved so far for  $1 \leq p' < \infty$ , but it can be verified that (3.8) holds for  $p' = \infty$ , too. We therefore have

$$\begin{aligned} \|V_c^{-T}\|_{p'} &= \max_{u \in \mathbb{R}^n} \frac{\|u\|_{p'}}{\|V_c^T u\|_{p'}} \geq \frac{\|v\|_{p'}}{\|V_c^T v\|_{p'}} \\ &\geq \min_{1 \leq k \leq k_{\max}} \left[ \frac{(n-k)!}{(n-1)!} \right]^2 \omega^{k-1} \times \frac{S_{n-1,p'}(\omega, \tau)}{n^{1/p'}}, \end{aligned}$$

as was to be shown. ■

**Theorem 3.1** *Assume (1.2) and (3.1). Then*

$$\kappa_p(V_c) \geq \min_{1 \leq k \leq k_{\max}} \left[ \frac{(n-k)!}{(n-1)!} \right]^2 \omega^{k-1} \times \max\{\ell^{1/p'}, \alpha^{n-1}\} \frac{S_{n-1,p'}(\omega, \tau)}{n^{1/p'}}. \quad (3.9)$$

Let  $\omega = \eta\alpha$ . Then

$$\kappa_p(V_c) \geq \min_{1 \leq k \leq k_{\max}} \left[ \frac{(n-k)!}{(n-1)!} \right]^2 \omega^{k-1} \times \frac{S_{n-1,1}(\eta, \tau)}{n^{1/p'}}. \quad (3.10)$$

*Proof:* Inequality (3.9) is a consequence of Lemmas 3.1 and 3.3. We now prove (3.10). Since  $a_{jn}(\omega, \tau) = \omega^{-j} a_{jn}(1, \tau)$  [10],  $S_{n-1,p'} = \left( \sum_{j=0}^{n-1} |\omega|^{-jp'} |a_{jn-1}(1, \tau)|^{p'} \right)^{1/p'}$ . Therefore it follows from (3.3) and (3.5) that

$$\begin{aligned} &n^{1/p'} \max_{1 \leq k \leq k_{\max}} \left[ \frac{(n-1)!}{(n-k)!} \right]^2 \frac{1}{\omega^{k-1}} \times \kappa_p(V_c) \\ &\geq \left( \sum_{j=0}^{n-1} \alpha^{jp} \right)^{1/p} \left( \sum_{j=0}^{n-1} |\omega|^{-jp'} |a_{jn-1}(1, \tau)|^{p'} \right)^{1/p'} \\ &\geq \sum_{j=0}^{n-1} |\eta|^{-j} |a_{jn-1}(1, \tau)| \\ &= S_{n-1,1}(\eta, \tau), \end{aligned}$$

by Hölder inequality. This yields (3.10). ■

For  $k_{\max} = 1$ , i.e.,  $\ell = n$  and  $k_1 = s = k_n = 1$  (and thus  $V_c = V$ ), (3.9) becomes one of the lower bounds for  $\kappa_p(V)$  in [10]. In general, we may also use (3.6), instead of (3.7), in estimating  $\|V_c^{-1}\|_p$ . Doing so will lead to a more complicated lower bound on  $\kappa_p(V_c)$ .

**Corollary 3.1** *Assume (1.2) and (3.1).*

$$\kappa_p(V_c) \geq \min_{1 \leq k \leq k_{\max}} \left[ \frac{(n-k)!}{(n-1)!} \right]^2 \alpha^{k-1} \times \frac{S_{n-1,1}(1, 0)}{n^{1/p'}}. \quad (3.11)$$

If, in addition, all  $\alpha_j \geq 0$ , then

$$\kappa_p(V_c) \geq \min_{1 \leq k \leq k_{\max}} \left[ \frac{(n-k)!}{(n-1)!} \right]^2 \left[ \frac{\alpha}{2} \right]^{k-1} \times \frac{S_{n-1,1}(1/2, 1)}{n^{1/p'}}. \quad (3.12)$$

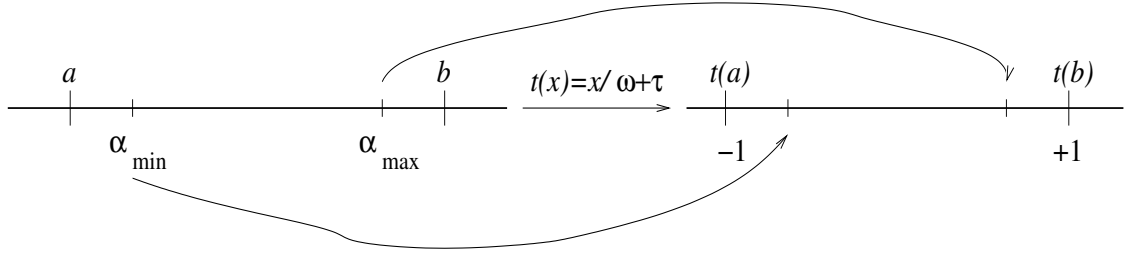
*Proof:* Apply (3.10) to  $[a, b] = [-\alpha, \alpha]$  (and thus  $\eta = 1$ ,  $\tau = 0$ , and  $\omega = \alpha$ ) to get (3.11). Apply (3.10) to  $[a, b] = [0, \alpha]$  (and thus  $\eta = 1/2$ ,  $\tau = -1$ , and  $\omega = \alpha/2$ ) to get (3.12). ■

Neither bounds in (3.11) and (3.12) are uniform, because both depend on  $\alpha$ . They do not yield useful lower bounds on  $\min_{\alpha_j} \kappa_p(V_c)$  or  $\min_{\alpha_j \geq 0} \kappa_p(V_c)$ . In fact, the minimums of both right-hand sides of (3.11) and (3.12) over either all  $\alpha_j \in \mathbb{R}$  or all  $\alpha_j \geq 0$  are zero! In the next two sections, we shall establish two uniform bounds using (3.9) for all real  $V_c$  and for those with nonnegative nodes.

**REMARK 3.1** Lemma 3.3 is made possible by Lemma 3.2 which was proved with the help of Markov's Inequality. Another classical inequality for the same purpose is Bernstein's Inequality [5, Page 233], using which we can obtain the following: *For  $0 \leq k \leq n$ , if  $a < \alpha_{\min} < \alpha_{\max} < b$ , then*

$$\left| \frac{d}{dx^k} T_n(x; \omega, \tau) \right| \leq \frac{n(n-1)s(n-k+1)}{\left[ \omega \sqrt{1 - \left( \frac{\max\{b - \alpha_{\max}, \alpha_{\min} - a\}}{\omega} \right)^2} \right]^k} \quad \text{for } x \in [a, b], \quad (3.13)$$

where  $\alpha_{\max} \equiv \max_j \alpha_j$  and  $\alpha_{\min} \equiv \min_j \alpha_j$ . This inequality improves (3.4) in the numerator part, but has complications in the denominator, and also it requires the interval  $[a, b]$  be (slightly) larger than the smallest interval containing all nodes, as follows.



This can be bad because larger  $[a, b]$  will weaken the effectiveness of  $S_{n,p'}(\omega, \tau)$  in the later bounds on  $\kappa_p(V_c)$ , for example  $S_{n,p'}(\omega, \tau)$  is decreasing in  $\omega$  [10].

## 4 Uniform bounds for all real $V_c$

We'll restrict ourselves to  $p = \infty$  because the availability of formulas for  $S_{n,1}$  in Section 2 that allow us to do analysis below. Equivalent relation (2.9) makes it possible to derive uniform lower bounds on  $\kappa_p(V_c)$  for  $p \neq \infty$ .

Let  $\Phi$  be the right-hand side of (3.9) with  $-a = b = \alpha$  for  $p = \infty$  (and thus  $p' = 1$ ):

$$\Phi = \min_{1 \leq k \leq k_{\max}} \left[ \frac{(n-k)!}{(n-1)!} \right]^2 \alpha^{k-1} \times \max\{\ell, \alpha^{n-1}\} \frac{S_{n-1,1}(\alpha, 0)}{n} \equiv \max\{\Phi_1, \Phi_2\},$$

where

$$\Phi_1 = \min_{1 \leq k \leq k_{\max}} \left[ \frac{(n-k)!}{(n-1)!} \right]^2 \alpha^{k-1} \times \frac{\ell S_{n-1,1}(\alpha, 0)}{n}, \quad (4.1)$$

$$\Phi_2 = \min_{1 \leq k \leq k_{\max}} \left[ \frac{(n-k)!}{(n-1)!} \right]^2 \alpha^{k-1} \times \frac{\alpha^{n-1} S_{n-1,1}(\alpha, 0)}{n}. \quad (4.2)$$

$\Phi$  is  $\Phi_1$  for  $\alpha \leq \ell^{1/(n-1)}$  and  $\Phi_2$  for  $\alpha \geq \ell^{1/(n-1)}$ .  $\Phi_2$  is increasing in  $\alpha$  for  $\alpha > 0$  because  $\alpha^{n-1}S_{n-1,1}(\alpha, 0)$  is increasing in  $\alpha$  for  $\alpha > 0$  [10]. Lemma 4.1 shows that  $\Phi_1$  is decreasing in  $\alpha$  for  $\alpha \leq 1$  for  $k_{\max}$  approximately no bigger than  $1 + (n-1)/\sqrt{2}$ . Between  $1 \leq \alpha \leq \ell^{1/(n-1)}$ ,  $\Phi = \Phi_1 \geq \Phi_2$  and  $\Phi_1 = \mathcal{O}_n(\Phi_2)$  because for  $1 \leq \alpha \leq \ell^{1/(n-1)}$ ,

$$\begin{aligned} \left[ \frac{(n - k_{\max})!}{(n-1)!} \right]^2 \frac{S_{n-1,1}(n^{1/(n-1)}, 0)}{n} &\leq \Phi_1 \leq \frac{n S_{n-1,1}(1, 0)}{n}, \\ \left[ \frac{(n - k_{\max})!}{(n-1)!} \right]^2 \frac{S_{n-1,1}(1, 0)}{n} &\leq \Phi_2 \leq \frac{n S_{n-1,1}(n^{1/(n-1)}, 0)}{n}, \end{aligned}$$

noting that  $\Phi_2$  is increasing in  $\alpha$  and that  $S_{n-1,1}(\alpha, 0)$  is decreasing in  $\alpha$ . These inequalities, together with,

$$\begin{aligned} n^{1/(n-1)} &= 1 + \frac{\ln(n-1)}{n-1} + \frac{2 + \ln^2(n-1)}{2(n-1)^2} + s, \\ S_{n-1,1}(n^{1/(n-1)}, 0) &\sim \frac{(1 + \sqrt{2})^{n-1}}{2(n-1)^{1/\sqrt{2}}} \end{aligned}$$

imply  $\Phi_1 = \mathcal{O}_n(\Phi_2)$  for  $1 \leq \alpha \leq \ell^{1/(n-1)}$ . Therefore we have the qualitative plot in Figure 1.1 for  $\Phi$ .

**Lemma 4.1** *Let  $k \geq 0$ .  $\alpha^k S_{n-1,1}(\alpha, 0)$  is decreasing in  $\alpha$  for  $\alpha \leq 1$  if*

$$k \leq \frac{n-1}{\sqrt{2}} \left[ 1 - (1 + \sqrt{2})^{-2n+2} \right] \sim \frac{n-1}{\sqrt{2}}. \quad (4.3)$$

*Proof:* We claim that under (4.3)  $\frac{d}{d\alpha} \alpha^k S_{n-1,1}(\alpha, 0) \leq 0$  for  $\alpha \leq 1$ . To this end, we notice

$$\frac{d}{d\alpha} \alpha^k S_{n-1,1}(\alpha, 0) = k \alpha^{k-1} S_{n-1,1}(\alpha, 0) + \alpha^k \frac{d}{d\alpha} S_{n-1,1}(\alpha, 0).$$

Now for  $\alpha \leq 1$ , by (2.2), we have

$$\begin{aligned} S_{n-1,1}(\alpha, 0) &\leq \frac{1}{2} \left[ \frac{1}{\alpha} + \sqrt{1 + \frac{1}{\alpha^2}} \right]^{n-1} [1 + \epsilon^{-2n+2}], \\ -\frac{d}{d\alpha} S_{n-1,1}(\alpha, 0) &\geq \frac{n-1}{2} \left[ \frac{1}{\alpha} + \sqrt{1 + \frac{1}{\alpha^2}} \right]^{n-2} [1 - \delta^{-2n+2}] \\ &\quad \times \left[ \frac{1}{\alpha^2} + \frac{1}{\alpha^2 \sqrt{1 + \alpha^2}} \right], \end{aligned}$$

where  $\epsilon = 1 + \sqrt{2}$  and  $\delta = 0$  for even  $n-1$ , and  $\epsilon = 0$  and  $\delta = 1 + \sqrt{2}$  for odd  $n-1$ . Therefore for  $\alpha \leq 1$

$$\begin{aligned} \frac{\frac{d}{d\alpha} \alpha^k S_{n-1,1}(\alpha, 0)}{(n-1)\alpha^{k-1} S_{n-1,1}(\alpha, 0)} &= \frac{k}{n-1} + \frac{\alpha \frac{d}{d\alpha} S_{n-1,1}(\alpha, 0)}{(n-1)S_{n-1,1}(\alpha, 0)} \\ &\leq \frac{k}{n-1} - \frac{\alpha \left[ \frac{1}{\alpha^2} + \frac{1}{\alpha^2 \sqrt{1 + \alpha^2}} \right]}{\frac{1}{\alpha} + \sqrt{1 + \frac{1}{\alpha^2}}} \frac{1 - \delta^{-2n+2}}{1 + \epsilon^{-2n+2}} \end{aligned}$$

$$\begin{aligned}
&= \frac{k}{n-1} - \frac{1}{\sqrt{1+\alpha^2}} \frac{1-\delta^{-2n+2}}{1+\epsilon^{-2n+2}} \\
&\leq \frac{k}{n-1} - \frac{1}{\sqrt{2}} \frac{1-\delta^{-2n+2}}{1+\epsilon^{-2n+2}} \\
&\leq \frac{k}{n-1} - \frac{1}{\sqrt{2}} \left[ 1 - (1+\sqrt{2})^{-2n+2} \right] \\
&\leq 0
\end{aligned}$$

upon using (4.3). ■

**Theorem 4.1** *If (4.3) with  $k = k_{\max} - 1$  holds, then*

$$\begin{aligned}
\kappa_{\infty}(V_c) &\geq \left[ \frac{(n-k_{\max})!}{(n-1)!} \right]^2 \frac{S_{n-1,1}(1,0)}{n} \\
&\sim \left[ \frac{(n-k_{\max})!}{(n-1)!} \right]^2 \frac{[1+\sqrt{2}]^{n-1}}{n}.
\end{aligned}$$

*Proof:* It can be verified that  $\Phi \geq \Phi_2|_{\alpha=1}$ . ■

## 5 Uniform bounds for $V_c$ with $\alpha_i \in [a, b]$ and $0 = a < b$

Let  $\Psi$  be the right-hand side of (3.9) with  $0 = a < b = \alpha$  for  $p = \infty$  (and thus  $p' = 1$ ):

$$\Psi = \min_{1 \leq k \leq k_{\max}} \left[ \frac{(n-k)!}{(n-1)!} \right]^2 \left[ \frac{\alpha}{2} \right]^{k-1} \times \max\{\ell, \alpha^{n-1}\} \frac{S_{n-1,1}(\alpha/2, 1)}{n} = \max\{\Psi_1, \Psi_2\},$$

where

$$\Psi_1 = \min_{1 \leq k \leq k_{\max}} \left[ \frac{(n-k)!}{(n-1)!} \right]^2 \left[ \frac{\alpha}{2} \right]^{k-1} \times \frac{\ell S_{n-1,1}(\alpha/2, 1)}{n}, \quad (5.1)$$

$$\Psi_2 = \min_{1 \leq k \leq k_{\max}} \left[ \frac{(n-k)!}{(n-1)!} \right]^2 \left[ \frac{\alpha}{2} \right]^{k-1} \times \frac{\alpha^{n-1} S_{n-1,1}(\alpha/2, 1)}{n}. \quad (5.2)$$

$\Psi$  is  $\Psi_1$  for  $\alpha \leq \ell^{1/(n-1)}$  and  $\Psi_2$  for  $\alpha \geq \ell^{1/(n-1)}$ .  $\Psi_2$  is increasing in  $\alpha$  for  $\alpha > 0$  because  $\alpha^{n-1} S_{n-1,1}(\alpha/2, 1)$  is increasing in  $\alpha$  for  $\alpha > 0$  [10]. Lemma 5.1 shows that  $\Psi_1$  is decreasing in  $\alpha$  for  $\alpha \leq 1$  for  $k_{\max}$  approximately no bigger than  $1 + (n-1)/\sqrt{2}$ . Between  $1 \leq \alpha \leq \ell^{1/(n-1)}$ ,  $\Psi = \Psi_1 \geq \Psi_2$  and  $\Psi_1 = \mathcal{O}_n(\Psi_2)$  because for  $1 \leq \alpha \leq \ell^{1/(n-1)}$ ,

$$\begin{aligned}
\left[ \frac{(n-k_{\max})!}{(n-1)!} \right]^2 \frac{1}{2^{k_{\max}-1}} \frac{S_{n-1,1}(n^{1/(n-1)}/2, 1)}{n} &\leq \Psi_1 \leq \frac{n S_{n-1,1}(1/2, 1)}{n}, \\
\left[ \frac{(n-k_{\max})!}{(n-1)!} \right]^2 \frac{1}{2^{k_{\max}-1}} \frac{S_{n-1,1}(1/2, 1)}{n} &\leq \Psi_2 \leq \frac{n S_{n-1,1}(n^{1/(n-1)}/2, 1)}{n},
\end{aligned}$$

noting that  $\Psi_2$  is increasing in  $\alpha$  and that  $S_{n-1,1}(\alpha/2, 1)$  is decreasing in  $\alpha$ . These inequalities, together with

$$\begin{aligned}
n^{1/(n-1)} &= 1 + \frac{\ln(n-1)}{n-1} + \frac{2 + \ln^2(n-1)}{2(n-1)^2} + s, \\
S_{n-1,1}(n^{1/(n-1)}/2, 1) &\sim \frac{(3 + 2\sqrt{2})^{n-1}}{2(n-1)^{1/\sqrt{2}}},
\end{aligned}$$

imply  $\Psi_1 = \mathcal{O}_n(\Psi_2)$  for  $1 \leq \alpha \leq \ell^{1/(n-1)}$ . Therefore we have the qualitative plot in Figure 1.1 for  $\Psi$ .

**Lemma 5.1** *Let  $k \geq 0$ .  $\alpha^k S_{n-1,1}(\alpha/2, 1)$  is decreasing in  $\alpha$  for  $\alpha \leq 1$  if*

$$k \leq \frac{n-1}{\sqrt{2}} \left[ 1 - (1 + \sqrt{2})^{-4(n-1)} \right]^{-1} \sim \frac{n-1}{\sqrt{2}}. \quad (5.3)$$

*Proof:* We claim that under (5.3)  $\frac{d}{d\alpha} \alpha^k S_{n-1,1}(\alpha/2, 1) \leq 0$  for  $\alpha \leq 1$ . To this end, we notice

$$\frac{d}{d\alpha} \alpha^k S_{n-1,1}(\alpha/2, 1) = k \alpha^{k-1} S_{n-1,1}(\alpha/2, 1) + \alpha^k \frac{d}{d\alpha} S_{n-1,1}(\alpha/2, 1).$$

Now for  $\alpha \leq 1$ , by (2.4), we have

$$\begin{aligned} S_{n-1,1}(\alpha/2, 1) &\leq \frac{1}{2} \left[ \frac{2}{\alpha} + 1 + \sqrt{\left(\frac{2}{\alpha} + 1\right)^2 - 1} \right]^{n-1} \left[ 1 + (3 + 2\sqrt{2})^{-2n+2} \right], \\ -\frac{d}{d\alpha} S_{n-1,1}(\alpha/2, 1) &\geq \frac{n-1}{2} \left[ \frac{2}{\alpha} + 1 + \sqrt{\left(\frac{2}{\alpha} + 1\right)^2 - 1} \right]^{n-2} \\ &\quad \times \frac{2}{\alpha^2} \left[ 1 + \frac{2+\alpha}{2\sqrt{1+\alpha}} \right]. \end{aligned}$$

Therefore for  $\alpha \leq 1$

$$\begin{aligned} \frac{\frac{d}{d\alpha} \alpha^k S_{n-1,1}(\alpha/2, 1)}{(n-1) \alpha^{k-1} S_{n-1,1}(\alpha/2, 1)} &= \frac{k}{n-1} + \frac{\alpha \frac{d}{d\alpha} S_{n-1,1}(\alpha/2, 1)}{(n-1) S_{n-1,1}(\alpha/2, 1)} \\ &\leq \frac{k}{n-1} - \frac{\alpha \frac{2}{\alpha^2} \left[ 1 + \frac{2+\alpha}{2\sqrt{1+\alpha}} \right]}{\frac{2}{\alpha} + 1 + \sqrt{\left(\frac{2}{\alpha} + 1\right)^2 - 1}} \frac{1}{1 + (3 + 2\sqrt{2})^{-2n+2}} \\ &= \frac{k}{n-1} - \frac{1}{\sqrt{1+\alpha}} \left[ 1 + (3 + 2\sqrt{2})^{-2n+2} \right]^{-1} \\ &\leq \frac{k}{n-1} - \frac{1}{\sqrt{2}} \left[ 1 + (3 + 2\sqrt{2})^{-2n+2} \right]^{-1} \\ &\leq 0 \end{aligned}$$

upon using (5.3). ■

**Theorem 5.1** *If (5.3) with  $k = k_{\max} - 1$  holds and all  $\alpha_i \geq 0$ , then*

$$\begin{aligned} \kappa_{\infty}(V_c) &\geq \left[ \frac{(n - k_{\max})!}{(n-1)!} \right]^2 \frac{1}{2^{k_{\max}-1}} \frac{S_{n-1,1}(1/2, 1)}{n} \\ &\sim \left[ \frac{(n - k_{\max})!}{(n-1)!} \right]^2 \frac{1}{2^{k_{\max}-1}} \frac{[1 + \sqrt{2}]^{2(n-1)}}{n}. \end{aligned}$$

*Proof:* It can be verified that  $\Psi \geq \Psi_2|_{\alpha=1}$ . ■

## 6 Conclusions

We have obtained several lower bounds on the condition number  $\kappa_p(V_c)$  of a real confluent Vandermonde matrix  $V_c$ . Two of them are uniform in the sense that they depend on  $n$ , the dimension of  $V_c$  only, while the others are either functions of  $n$  and  $\alpha$  or  $n$  and the interval  $[a, b]$  that contains all  $\alpha_j$ . These bounds grow exponentially for any fixed  $k_{\max}$ , much as expected. Qualitative behaviors of our general lower bound (3.9) for  $-a = b = \alpha$  and for  $0 = a < b = \alpha$  are plotted in Figure 1.1. While it is not clear in general if (any of) our bounds are asymptotically optimal, in contrast to those for Vandermonde matrices by Beckermann [2] and recently by the author [10], our bounds are unlikely to be asymptotically optimal if  $k_{\max}$  also grows, e.g., linearly in  $n$ . This is illustrated by the extreme example  $k_{\max} = n$ , as we commented in Section 1.

We have focused on real confluent Vandermonde matrices here. It is conceivable that there would be much better conditioned complex confluent Vandermonde matrices or confluent Vandermonde-like matrices. This is partly an intuition one might get from that although real Vandermonde matrices are very ill-conditioned [7, 2, 10, 12], there exist very well-conditioned complex Vandermonde matrices and Vandermonde-like matrices [6, 11]. We plan to investigate this issue in future work.

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