

# Minimal Curves

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- What is the original algorithm?
- An example
- How is the new algorithm an improvement?

# Some Notation

- We will be working in  $\mathbb{P}^3$ .
- $K$  will denote an algebraically closed field.
- $R = K[X, Y, Z, T]$  will be a polynomial ring.
- We also note that a curve  $C \subseteq \mathbb{P}^3$  is a closed subscheme of dimension 1 that is taken to be locally Cohen-Macaulay (in particular, no embedded points).
- Let  $C \subset \mathbb{P}^3$  be a curve. Then the Hartshorne-Rao module of  $C$  is:

$$M(C) = \mathbf{H}_m^1 \left( R/I_C \right)$$

(f.g. since the curve is locally Cohen-Macaulay) .

# What is a minimal curve?

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Proof: We will prove that there is a minimal left shift of  $M$ .

Let  $M$  be a f.g. unmixed module associated to a curve.

Then our claim is that  $H_m^1(R/I_C) \cong M(e)$  is not possible for  $e \gg 0$ .

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Then our claim is that  $H_m^1(R/I_C) \cong M(e)$  is not possible for  $e \gg 0$ . We will proceed by contradiction.

# What is a minimal curve?

Assume that  $H_m^1(R/I_C) \cong M(e) \neq 0$  such that

$$\left[ H_m^1(R/I_C) \right]_j = [M]_{e+j} = 0, \quad \text{for all } j \geq 0$$

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Take  $\ell \in [R]_1$  general. Then we have the following exact sequence:

$$0 \longrightarrow A(-1) \xrightarrow{\ell} A \longrightarrow A/\ell A \longrightarrow 0$$

where  $A = R/I_C$ .

# What is a minimal curve?

This induces the long exact cohomology sequence:

$$\dots \longrightarrow H_m^0(A) \longrightarrow H_m^0(A/\ell_A) \longrightarrow H_m^1(A)(-1) \xrightarrow{\ell} H_m^1(A) \longrightarrow \dots$$

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But,  $H_m^0(A) = 0$ . Hence, we have:

$$0 \longrightarrow H_m^0(A/\ell_A) \longrightarrow M(e-1) \longrightarrow M(e)$$

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since  $H_m^1(A) \cong M(e)$ . Now, we look in degree  $a+1$ :

$$0 \longrightarrow \left[ H_m^0(A/\ell_A) \right]_{a+1} \longrightarrow [M(e-1)]_{a+1} \longrightarrow [M(e)]_{a+1}$$

where  $a = \max \{ j \mid [M(e)]_j \neq 0 \}$ .

# What is a minimal curve?

Then  $a < 0$  by the choice of  $e$ . Moreover,  $[M(e - 1)]_{a+1} \neq 0$  and  $[M(e)]_{a+1} = 0$ .

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But,  $\left[ H_m^0 \left( A/\ell A \right) \right]_{a-1} \neq 0$  where  $a < 0$ . A contradiction.

Therefore, there is a left-most shift. *QED*

# What is a minimal curve?

- Definition of a Minimal Curve: Let  $\mathcal{L}$  be a liaison class. A curve  $C \in \mathcal{L}$  whose Hartshorn-Rao module corresponds to the left-most possible shift is called a minimal curve.

# Why are minimal curves important?

- Lazarsfeld-Rao Property: Let  $\mathcal{L}$  be a biliaison class of curves in  $\mathbb{P}^3$ , then we say that  $\mathcal{L}$  has the Lazarsfeld-Rao Property if the following conditions hold:
  1. If  $C_1$  and  $C_2$  are minimal curves, then there is a deformation from one to the other through subschemes which are all minimal curves.
  2. Let  $C_0$  be a minimal curve. Given  $C \in \mathcal{L}$ , then there exists a sequence of subschemes  $C_0, C_1, \dots, C_t$  such that for all  $i$ ,  $1 \leq i \leq t$ ,  $C_i$  is a basic double link of  $C_{i-1}$ , and  $C$  is a deformation of  $C_t$ .

# Why are minimal curves important?

- Theorem: Every liaison class of non-arithmetically Cohen-Macaulay curves of  $\mathbb{P}^3$  has the Lazarsfeld-Rao Property.

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- Theorem: Every biliaison class of non-arithmetically Cohen-Macaulay curves of  $\mathbb{P}^3$  has the Lazarsfeld-Rao Property.
- So, if we were able to find a minimal curve in the biliaison class of a given module, we could recover any other curve in the biliaison class by a sequence of basic double links and possibly a deformation.

# The original algorithm

Let  $M$  be the Hartshorne-Rao Module of the curve  $C$ .







# The original algorithm

We note that the exact sequence

$$0 \longrightarrow I \longrightarrow R \longrightarrow R/I \longrightarrow 0$$

induces the long exact cohomology sequence:

$$\dots \longrightarrow H_m^1(R) \longrightarrow H_m^1(R/I) \longrightarrow H_m^2(I) \longrightarrow H_m^2(R) \longrightarrow \dots$$

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But,  $R$  is free, so  $H_m^1(R) = H_m^2(R) = 0$ . Hence, we get the isomorphism

$$H_m^1(R/I) \cong H_m^2(I)$$

# The original algorithm

Using this same idea, we can get the following isomorphisms:

$$H_m^1(R/I) \cong H_m^2(I) \cong H_m^2(L)(-s) \cong H_m^1(Y)(-s) \cong M(-s)$$

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However, for  $I$  to define a curve, we need to check that  $I$  is saturated. We know that  $I$  is saturated iff  $H_m^0(R/I) = 0$ . But

$$H_m^0(R/I) \cong H_m^1(I) \cong H_m^1(L)(-s) = 0$$

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We have the following exact sequence:

$$0 \longrightarrow F \xrightarrow{\varphi} L \longrightarrow I(s) \longrightarrow 0$$

where  $F = \bigoplus_{i=1}^{r-1} R(-d_i)$  with  $d_1 \leq d_2 \leq \dots \leq d_{r-1}$ .

# The original algorithm

Since we have free resolutions for  $F$  and  $L$ , we can use the mapping cone to get a free resolution for  $I(s)$  (not necessarily minimal).

$$\begin{array}{ccccccc}
 0 & \longrightarrow & F & \xrightarrow{\varphi} & L & \longrightarrow & I(s) \longrightarrow 0 \\
 & & \uparrow & & \uparrow & & \uparrow \\
 & & F & & F_2 & & F_2 \\
 & & \uparrow & & \uparrow & & \uparrow \\
 & & 0 & & F_3 & & F_3 \oplus F \\
 & & & & \uparrow & & \uparrow \\
 & & & & F_4 & & F_4 \\
 & & & & \uparrow & & \uparrow \\
 & & & & 0 & \cdot & 0 \cdot \cdot \cdot
 \end{array}$$

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If we considered  $\sigma_2$  as a matrix, then the columns are the minimal generators of  $L$ .

Then the matrix defining  $\varphi$  consists of  $r-1$  of these columns where we take into consideration the degrees among the generators, i.e. the  $d_i$ 's.

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Moreover,  $\text{coker}(\varphi) \cong I$  iff  $\text{ht}(I_{r-1}) = 2$ .

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Moreover,  $\text{coker}(\varphi) \cong I$  iff  $\text{ht}(I_{r-1}) = 2$ . It ends up that the saturation of the ideal generated by the  $(r-1) \times (r-1)$  minors of the matrix corresponding to  $\varphi$  is the defining ideal of the minimal curve.

# Improvements of the new algorithm

- The newer algorithm is done in  $K[u]$  which is a much simpler polynomial ring.

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- The newer algorithm is done in  $K[u]$  which is a much simpler polynomial ring.
- In  $K[u]$ , the problem is reduced to applying the Smith Normal Form algorithm to the new matrix.

# References

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