Intro to Contemporary Math Dijkstra's Algorithm

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Agenda

- ▶ Dijkstra's Algorithm
 - Procedure
 - Evaluation

Distance Between Vertices

- ▶ Definition: A **distance** between two vertices is the total weight of a path between them
- ► Goal: Want shortest path (path with smallest total weight) which gives shortest distance between a starting vertex and destination vertex

Dijkstra's Algorithm Overview, Steps 1-2

Instead of testing entire paths, we build one.

Step 1: Mark the destination vertex as **current**. Record its **distance** as **zero**: its has distance zero from the destination vertex.

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Dijkstra's Algorithm Overview, Steps 1-2

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Step 2: For each unvisited vertex **connected** to the current one, add these numbers:

- ▶ the weight of the edge from this vertex to current one
- the distance from the current vertex to the destination

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Step 2a: Record the sum (distance) and the current vertex next to this vertex's label UNLESS the vertex is already recorded with a smaller distance.

Dijkstra's Algorithm Steps 2-4

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Step 2a: Record the sum (distance) and the current vertex next to this vertex's label UNLESS the vertex is already recorded with a smaller distance.

Step 3: Mark the **current** vertex as **visited**.

Step 4: Mark the unvisited vertex with the smallest distance as current, and repeat from step 2.

Dijkstra's Algorithm Steps 2a-4+End

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Step 3: Mark the **current** vertex as **visited**.

Step 4: Mark the unvisited vertex with the smallest distance as current, and repeat from step 2.

Stop when all vertices (except the start vertex) are visited.

Dijkstra's Algorithm Evaluation

Optimal: Guaranteed to find shortest path (path of least weight)

Algorithm must have safeguards to avoid missing shortest path

- ► Step 4 prevents algorithm from missing shorter paths from an unvisited vertex to the destination.
- Unvisited vertices with higher number records may get rerecorded in Step 2a. Going to them first will miss shorter paths from them!

Dijkstra's Algorithm Evaluation

Efficient: Computers can use it to find a shortest path quickly Algorithm must have ways to avoid testing bad paths

- ► Step 3 prevents the algorithm from going in circles.
- ► A vertex is visited when we know the shortest distance from it to the destination, so the algorithm knows not to try other paths from it.

Dijkstra's vs Brute Force

- ► Both are optimal in theory they find the shortest path
- ▶ Dijkstra's is efficient, but
- Brute Force is not efficient because it tests every path, including bad ones (that go in circles, make U-turns, etc.)

Dijkstra's vs Brute Force

For comparison, on a graph with 25 vertices,

- Dijkstra's takes 625 calculations
- ► Brute Force takes 10²⁵ calculations that's 1 with 25 zeroes after it.

Next Time

► Dijkstra's Algorithm: Interpreting Output

Bibliography

► Lippman, David. Math in Society. 2nd ed. 16 November 2013. http://www.opentextbookstore.com/mathinsociety/current2.php?chapter=GraphTheory.pdf. Web.