## Some Corrections to the Foundations of Combinatorial Optimization

- 1. p.7, l.5, change " $a'_{36} \neq 0$ " to " $a'_{37} \neq 0$ ".
- 2. p.11, in Theorem (Greedy characterization of matroids), define a base of an independence system to be a maximal independent set.
- 3. p.17, l.12, change " $\mathbf{R}^{2^{E(G)}}$ " to " $\mathbf{R}^{2^{E(M)}}$ "
- 4. p.21, l.10, "notice that J + e contains a unique circuit".
- 5. p.21, last line, "by noting that  $\alpha = 0$  would imply  $\mathcal{F}(T) = \mathcal{P}_{\mathcal{I}(M)}$ , and  $\alpha < 0$  would imply  $T = \emptyset$ ."
- 6. p.24, Bellman-Ford Algorithm, step 2, I think you must state " $f_k(w) := \min(\{f_{k-1}(w)\} \cup \{f_{k-1}(t(e)) + c(e) : e \in \delta_G^-(w)\})$ ".
- 7. p.26, l.5 from the bottom, "At any stage of the algorithm".
- 8. p.28, try to clarify the statement, "It cannot be that  $w^*$  must be used by such an F' before the last interior vertex j, because (a) implies that there is a minimum-weight v-j dipath that does not use  $w^*$ ."
- 9. p.29, Knapsack Problem/Exercise. I suggest in (a) giving a knapsack problem in which the slack variable is necessarily positive in the optimal solution.
- 10. p.33, l.-10, "that is, those induced by the rigid motions of  $\mathbf{R}^{d}$ ".
- 11. p.35, l.6, "for any pair of matroids  $M_1, M_2$ ".
- 12. p.46, in the Exercise, " $\mathcal{P}_{\mathcal{I}(M_1)} \cap \mathcal{P}_{\mathcal{I}(M_2)} \cap \mathcal{P}_{\mathcal{I}(M_3)}$ ".
- 13. p.49, line before Section 4.1, "minimum-weight matchings".
- 14. p.50, l.-10. I think this should say "there is some (alternating) path of G.C with more vertices in Y than in X".
- 15. p.52, in Claim 1, mention that  $S^1|_{G_1} \cup S^2|_{G_2}$  is a maximizing matching for  $\alpha$ .
- 16. p.62, in the definition of Eulerian tour, specify that the directed walk must be closed.
- 17. p.67, l.3, "optimization".

- 18. p.73, l.19, "therefore  $\sigma_i^k \leq \sigma_i^{k+1}$ ".
- 19. p.73, l.-5, "Let's".
- 20. p.79, in the Example, "The choice of  $u_1 = 0$ ,  $u_2 = 1/2$ ,  $u_3 = 1/2$  yields the cutting plane  $-5x_1 3x_2 \le -13$ ."
- 21. p.85, equation ( $\overline{2}$ ), change " $\frac{15}{22}$ " to " $\frac{7}{22}$ ".
- 22. p.93, l.1, "Assume that the data for IP".
- 23. p.94, 1.2, "is  $z^*$ ."
- 24. p.94, l.-6, "is not too large".
- 25. p.95, first line of Proof, "is a feasible solution".
- 26. p.95, fourth line of Proof, "are satisfied by  $x^*$ ".
- 27. p.98, l.4, "lifting the coefficients of  $x_1$  and  $x_5$ " (double check this).
- 28. p.100, l.2, "with respect to the".
- 29. p.106, Problem, l-2, change "more than" to "at least". (This only makes a difference in the case n = 1.)
- 30. p.109, reverse the arc with cost 5/6 to go from 5 to 0.
- 31. p.109, l=8, "from vertex 4 to vertex 1".
- 32. p.110, fourth line before the Exercise, "potential".
- p.110, just before the Exercise, clarify the branching procedure for the knapsack problem.
- 34. p.113, definition of f(S), I think this should be " $f(S) := r_M(S) \sum_{e \in S} x_e^*$ ."
- 35. p.114, 1–7, " $f'(x) := \sum_{j=1}^{m} \lambda_j f(S(x^j))$ ".
- 36. p.115, l.2, "coefficients. Therefore".
- 37. p.115, you need a summation on the right-hand side of both inequalities (\*  $\leq$ ) and (\*  $\leq$ ).
- 38. p.127, l.-3, " $\sum_{i=1}^{m} y_i a_{ij} = 0$ , for j = 1, 2, ..., n".